



OpenCores.Org

# R65C02 IP Core Specification

*Author: Jens Gutschmidt*  
*opencores@vivare-services.com*

**Rev. 1.01**  
**February 19, 2021**



*This page has been intentionally left blank.*

## Revision History

Rev	Date	Author	Description
0.1	12/18/06	Jens Gutschmidt	First Draft
0.2	01/02/07	Jens Gutschmidt	Pictures of FSM's
0.3	08/20/08	Jens Gutschmidt	<ul style="list-style-type: none"> <li>- Tables and timing diagrams</li> <li>- Deleting pictures of FSM</li> </ul>
0.4	10/01/08	Jens Gutschmidt	<ul style="list-style-type: none"> <li>- New ideas for timing diagrams</li> </ul>
0.5	01/02/09	Jens Gutschmidt	<ul style="list-style-type: none"> <li>- Textual changes / spell checking</li> <li>- Insert R6502_TC block diagram</li> </ul>
0.6	02/01/09	Jens Gutschmidt	<ul style="list-style-type: none"> <li>- Work on Timing Diagrams</li> </ul>
0.7	23/02/09	Jens Gutschmidt	<ul style="list-style-type: none"> <li>- Added new instruction codes (Timing diagrams for all new op codes missing)</li> </ul>
0.8	16/04/14	Jens Gutschmidt	<ul style="list-style-type: none"> <li>- Corrected description for BIT</li> <li>- Corrected timing for JSR, RTS</li> </ul>
0.9	09-Sep-2018	Jens Gutschmidt	<ul style="list-style-type: none"> <li>- Corrected ADC/SBC descriptions (D=1 -&gt; 1 additional cycle)</li> <li>- Corrected all branch descriptions (page crossing computed wrong)</li> <li>- PHP writes always E&amp;B = '1' to stack</li> <li>- Adding Interrupt section</li> <li>- Adding ADC / SBC Decimal Mode</li> </ul>
0.91	04-Oct-2018	Jens Gutschmidt	<ul style="list-style-type: none"> <li>- Corrected the instruction table (all ABS,X ABS,Y and (IND),Y address modes are depend on page crossing – wrong in vendor's documentation)</li> </ul>
1.00	14-Oct-2018	Jens Gutschmidt	<ul style="list-style-type: none"> <li>- Section "Timing Diagrams" has been deleted</li> </ul>
1.01	18-Feb-2021	Jens Gutschmidt	<ul style="list-style-type: none"> <li>- Two core versions available now (1.53/"NORMAL"-2.00/"HIGH SPEED")</li> </ul>

# Contents

<b>1</b>	<b>INTRODUCTION .....</b>	<b>1</b>
<b>2</b>	<b>ARCHITECTURE .....</b>	<b>2</b>
<b>3</b>	<b>OPERATION .....</b>	<b>6</b>
	<i>ADC</i> .....	7
	<i>AND</i> .....	8
	<i>ASL</i> .....	9
	<i>BBR</i> .....	10
	<i>BBS</i> .....	11
	<i>BCC</i> .....	12
	<i>BCS</i> .....	13
	<i>BEQ</i> .....	14
	<i>BIT</i> .....	15
	<i>BMI</i> .....	16
	<i>BNE</i> .....	17
	<i>BPL</i> .....	18
	<i>BRA</i> .....	19
	<i>BRK</i> .....	19
	<i>BVC</i> .....	20
	<i>BVS</i> .....	21
	<i>CLC</i> .....	22
	<i>CLD</i> .....	22
	<i>CLI</i> .....	23
	<i>CLV</i> .....	23
	<i>CMP</i> .....	24
	<i>CPX</i> .....	24
	<i>CPY</i> .....	25
	<i>DEC</i> .....	25
	<i>DEX</i> .....	26
	<i>DEY</i> .....	26
	<i>EOR</i> .....	27
	<i>INC</i> .....	28
	<i>INX</i> .....	28
	<i>INY</i> .....	29
	<i>JMP</i> .....	29
	<i>JSR</i> .....	30
	<i>LDA</i> .....	30
	<i>LDX</i> .....	31
	<i>LDY</i> .....	31
	<i>LSR</i> .....	32
	<i>NOP</i> .....	32
	<i>ORA</i> .....	33
	<i>PHA</i> .....	33
	<i>PHP</i> .....	34
	<i>PHX</i> .....	34
	<i>PHY</i> .....	34
	<i>PLA</i> .....	34
	<i>PLP</i> .....	35
	<i>PLX</i> .....	35

---

<i>PLY</i> .....	35
<i>ROL</i> .....	36
<i>ROR</i> .....	36
<i>RMB</i> .....	37
<i>RTI</i> .....	37
<i>RTS</i> .....	37
<i>SBC</i> .....	38
<i>SEC</i> .....	38
<i>SED</i> .....	39
<i>SEI</i> .....	39
<i>SMB</i> .....	40
<i>STA</i> .....	40
<i>STX</i> .....	41
<i>STY</i> .....	41
<i>STZ</i> .....	41
<i>TAX</i> .....	42
<i>TAY</i> .....	42
<i>TRB</i> .....	43
<i>TSB</i> .....	44
<i>TSX</i> .....	44
<i>TXA</i> .....	45
<i>TXS</i> .....	45
<i>TYA</i> .....	46
<b>4</b> .....	<b>47</b>
<i>Registers</i> .....	47
<i>List of Registers</i> .....	47
<i>PSW – Description</i> .....	47
<b>5</b> .....	<b>48</b>
<i>Clocks</i> .....	48
<b>6</b> .....	<b>49</b>
<i>Interrupts</i> .....	49
<i>NMI</i> .....	49
<i>IRQ</i> .....	50
<i>BRK</i> .....	50
<b>APPENDIX A</b> .....	<b>51</b>
<b>ADC / SBC DECIMAL MODE</b> .....	<b>51</b>
<b>INSTRUCTION TABLE</b> .....	<b>53</b>
<b>INDEX</b> .....	<b>54</b>
<b>LIST OF FIGURES</b>	
Figure (1): R65C02 architecture.....	2
Figure (2): R65C02_TC IP core architecture .....	5
Figure (3): Interrupt NMI – Timing Diagram .....	49
Figure (4): Interrupt IRQ – Timing Diagram .....	50
<b>LIST OF TABLES</b>	

Table (1): ADC – Short Reference .....	7
Table (2): AND – Short Reference .....	8
Table (3): ASL – Short Reference .....	9
Table (4): BBR – Short Reference .....	10
Table (5): BBS – Short Reference .....	11
Table (6): BCC – Short Reference .....	12
Table (7): BCS – Short Reference .....	13
Table (8): BEQ – Short Reference .....	14
Table (9): BIT – Short Reference .....	15
Table (10): BMI – Short Reference .....	16
Table (11): BNE – Short Reference .....	17
Table (12): BPL – Short Reference .....	18
Table (13): BRA – Short Reference .....	19
Table (14): BRK – Short Reference .....	19
Table (15): BVC – Short Reference .....	20
Table (16): BVS – Short Reference .....	21
Table (17): CLC – Short Reference .....	22
Table (18): CLD – Short Reference .....	22
Table (19): CLI – Short Reference .....	23
Table (20): CLV – Short Reference .....	23
Table (21): CMP – Short Reference .....	24
Table (22): CPX – Short Reference .....	24
Table (23): CPY – Short Reference .....	25
Table (24): DEC – Short Reference .....	25
Table (25): DEX – Short Reference .....	26
Table (26): DEX – Short Reference .....	26
Table (27): EOR – Short Reference .....	27
Table (28): INC – Short Reference .....	28
Table (29): INX – Short Reference .....	28
Table (30): INY – Short Reference .....	29
Table (31): JMP – Short Reference .....	29
Table (32): JSR – Short Reference .....	30
Table (33): LDA – Short Reference .....	30

---

---

Table (34): LDX – Short Reference .....	31
Table (35): LDY – Short Reference .....	31
Table (36): LSR – Short Reference .....	32
Table (37): NOP – Short Reference .....	32
Table (38): OR – Short Reference .....	33
Table (39): PHA – Short Reference .....	33
Table (40): PHP – Short Reference .....	34
Table (41): PHX – Short Reference .....	34
Table (42): PHY – Short Reference .....	34
Table (43): PLA – Short Reference .....	34
Table (44): PLP – Short Reference .....	35
Table (45): PLX – Short Reference .....	35
Table (46): PLY – Short Reference .....	35
Table (47): ROL – Short Reference .....	36
Table (48): ROR – Short Reference .....	36
Table (49): RMB – Short Reference .....	37
Table (50): RTI – Short Reference .....	37
Table (51): RTS – Short Reference .....	37
Table (52): SBC – Short Reference .....	38
Table (53): SEC – Short Reference .....	38
Table (54): CLD – Short Reference .....	39
Table (55): SEI – Short Reference .....	39
Table (56): SMB – Short Reference .....	40
Table (57): STA – Short Reference .....	40
Table (58): STX – Short Reference .....	41
Table (59): STY – Short Reference .....	41
Table (60): STZ – Short Reference .....	41
Table (61): TAX – Short Reference .....	42
Table (62): TAY – Short Reference .....	42
Table (63): TRB – Short Reference .....	43
Table (64): TSB – Short Reference .....	44
Table (65): TSX – Short Reference .....	44
Table (66): TXA – Short Reference .....	45

---



Table (67): TXS – Short Reference .....	45
Table (68): TYA – Short Reference .....	46

---

# 1 Introduction

---

The Central Processing Unit (CPU) 6502 was introduced at 1976 by Commodore Computers. It is an 8 Bit processor which is well known worldwide. The also most known computer system in the 70s/80s was the APPLE based on this famous 6502.

In this century building of little computer systems based on an 8 Bit architecture is easier than ever before. Many CAD/CAM tool are existing on the market to give you help to do this job.

In the last decade the technology give us more and more possibilities to reach our dreams – higher, faster, wider. FPGA's (Field Programmable Gate Arrays) and CPLD's (Complex Programmable Logic Developes) are coming up. They shorten the time of development, simulation and verification of such systems dramatically. On the other hand stands the rising complex of designs and well knowing of desired description languages – VHDL (Very high scale integration Hardware Description Language) or Verilog.

Many operations and functions of computer tasks aren't need really the power of “modern” processors today like Intel's Pentium. Nevertheless it is not usable – sometimes also impossible - to build simple systems with discrete IC's. The necessity of flexibility compel us the using of high tech parts and tools...to build a “simple” system.

The using of standard IP's (Intellectual Properties) is the key to reach that goal.

Now, talking about the specification of the **R65C02 True Cycle Core IP**. Written in VHDL but developed with Mentor's HDL Designer. The R65C02 is very easy to handle for people were have experience with other cpu. Also suitable for beginners who will learning to build her own first cpu system.

# 2 Architecture

The following figure shows the internal architecture of the Rockwell's R65C02.

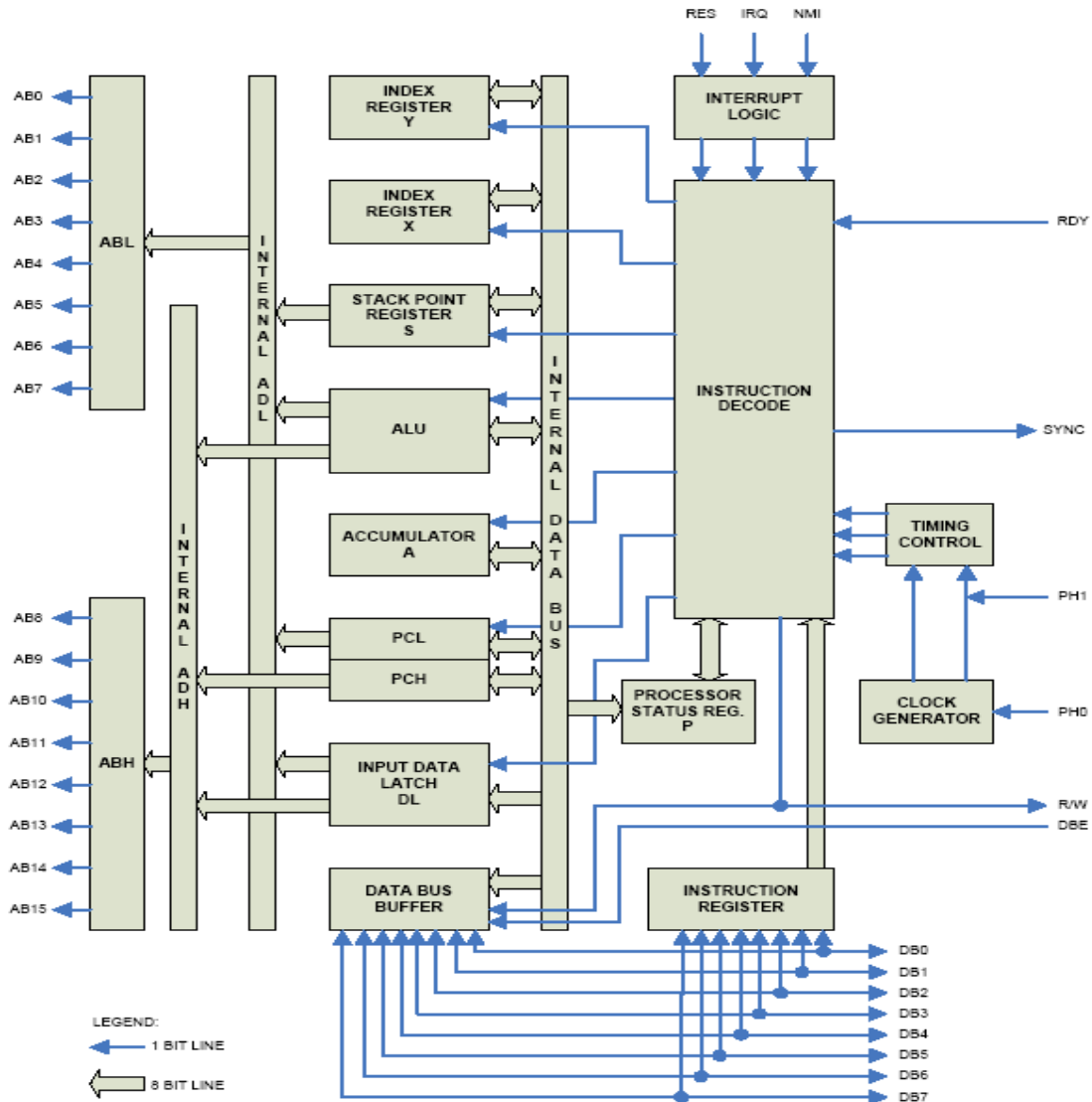


Figure (1): R65C02 architecture

First, for more clearance and better understanding of the **R65C02 True Cycle Core IP** let me summarize the most important quality of Rockwell's R65C02.

This architecture based on asynch logic. The two phase clock is a special attribute of that. The Data Bus is only active while clock phase PH1 is "1" for reading and writing (PH0 is "0"). All transfers of data are occurring at active phase of PH1.

The whole R65C02 is fully statically, so all registers holding her values if the clocking will be going to D.C. Internally all operation of 16 Bit are splitted into two 8 Bit busses. This is very interesting and causes one more clock cycle in some operations (see “PCL”, “PCH” -> Program Counter Low/High) which can operate over page boundaries.

Address lines are 16 Bit wide. 65.536 addresses for byte access are possible. There are no differences exist between memory and I/O cycles. All I/O devices must be connected via memory mapped I/O. The mnemonics are also all the same for both. The usage of address space is normally organized from top to the button “ROM – I/O – RAM” (0xFFFF – 0x0000).

Only one pin  $\overline{R/W}$  controls the read and writes operations.

The R65C02 can hold in every cycle by applying RDY. Wait states can be generated by using RDY.

There are two interrupts  $\overline{IRQ}$  and  $\overline{NMI}$ .  $\overline{IRQ}$  is mask able,  $\overline{NMI}$  is not. Note that  $\overline{RES}$  is internally handled as an interrupt, but is resetting the R65C02 to the starting point of operation at vector address 0xFFFF. Every interrupt has its own vector which can point elsewhere to the 16 Bit address space. The vectors are located from 0xFFFF to 0xFFFFA.

The Stack Pointer is 256 Bytes deep and is hard-wired to 0x01FF (internally 0x0200 –0x01FF going to the address bus) for the first memory location and grows downward to 0x0100. Decrementing again will produce a wraparound back to 0x01FF.

Arithmetical operations like ADC and SBC can handle binary values from 0-255 and decimal values from 0-99 without using other op codes. Only one user changeable bit into the Status Register determine the exact arithmetic mode of that operations.

The R65C02 generally operates into a pipeline mode. That means that a finish of an OP code falls every time into a fetch cycle of the next OP. This save one clock cycle

Every operation consume between one and eight clock cycles.

To build a useful and backward software and hardware compatible VHDL Core for the R65C02 many unavoidable changes may be forced to simulation and verification before any real core will be made.

The requirements:

- Vendor independent for implementation in any FPGA
- True Cycle for all operations as described in original publications  
*Because the original publications until late 80's shows wrong cycle counts for some R65C02 instructions, the r65c02\_tc's cycle timing was finally referenced against [4] (Programming the 65816).*
- No "Mixed Mode" – Only R65C02, not R6502 -> for performance and area
- No fantastic or useful extensions – as like the original R65C02 as possible
- Only one clock source
- Full synch design
- Operating speed  $f_{max} > 100\text{MHz}$  for most public used FPGA devices
- Easily to change for future requirement – building variants
- The activity on data and address busses since the execution of OP codes is not forced to be like the original – may be or maybe not.

The design based on finite state machines (fsm). Some registers are for internally use only and stores temporally values. These registers are not accessible by the programmer.

The next picture shows the hierarchical structure of the R65C02 IP core.

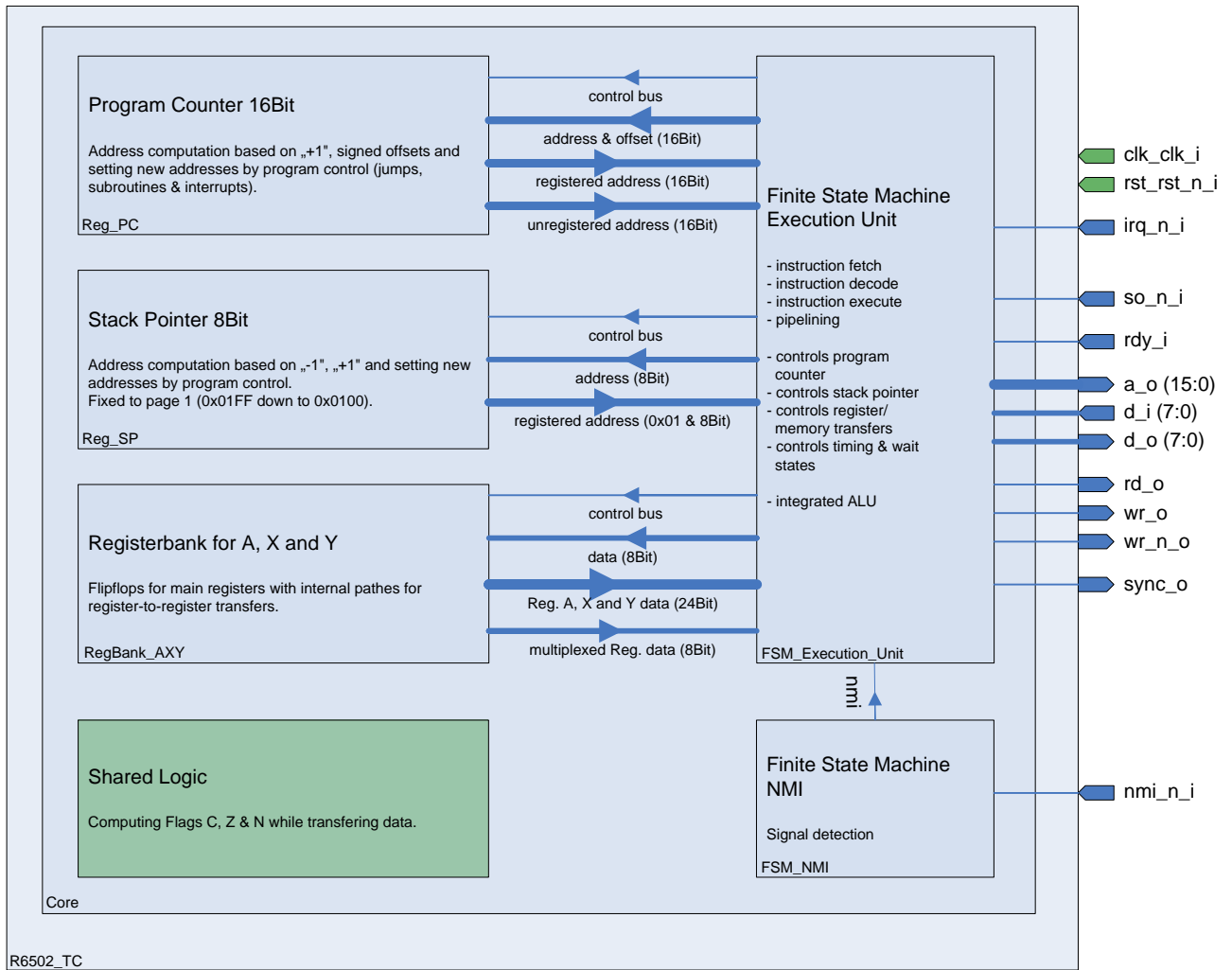


Figure (2): R65C02\_TC IP core architecture

The current IP core was divided into two versions for now on.

One still exists as traditional coding style and the other one introduce an optimizing technique for performance improvements.

Starting with v1.53 this branch left in traditional coding style.

Starting with v2.00 this branch contains the optimizing technique for speed.

If there are requirements to reduce the area of the core please use v1.xx instead of v2.xx.

# 3 Operation

---

# ADC

Operation: Add memory or immediate value to accumulator A with carry ( $A + M + C \rightarrow A, C$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Immediate	ADC #Oper	69	2	2+	√	√	-	-	-	-	√	√
Zero Page	ADC Oper	65	2	3+	√	√	-	-	-	-	√	√
Zero Page, X	ADC Oper, X	75	2	4+	√	√	-	-	-	-	√	√
Absolute	ADC Oper	6D	3	4+	√	√	-	-	-	-	√	√
Absolute, X	ADC Oper, X	7D	3	4+*	√	√	-	-	-	-	√	√
Absolute, Y	ADC Oper, Y	79	3	4+*	√	√	-	-	-	-	√	√
(Indirect, X)	ADC (Oper, X)	61	2	6+	√	√	-	-	-	-	√	√
(Indirect), Y	ADC (Oper), Y	71	2	5+*	√	√	-	-	-	-	√	√
(Indirect) <i>NEW</i>	ADC (Oper)	72	2	5+	√	√	-	-	-	-	√	√

\* => Add 1 if page boundary is crossed

+ => Add 1 if mode is "DECIMAL" (D='1')

Table (1): ADC – Short Reference

Example Immediate (assumed A=\$FE, C=0, D=0):

Address	Bytes	Mnemonic	
\$9000	69 03	ADC # \$03	;
\$9002	...	next OP	; A is now \$01, C=1, N=0, Z=0, V=0

1111 1110 (\$FE)  
ADC 0000 0011 (\$03)  
0000 0001 (\$01), C=1, N=0, Z=0, V=0



# AND

Operation: “AND” memory or immediate value with accumulator A ( $A \cap M \rightarrow A$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Immediate	AND #Oper	29	2	2	√	-	-	-	-	-	√	-
Zero Page	AND Oper	25	2	3	√	-	-	-	-	-	√	-
Zero Page, X	AND Oper, X	35	2	4	√	-	-	-	-	-	√	-
Absolute	AND Oper	2D	3	4	√	-	-	-	-	-	√	-
Absolute, X	AND Oper, X	3D	3	4*	√	-	-	-	-	-	√	-
Absolute, Y	AND Oper, Y	39	3	4*	√	-	-	-	-	-	√	-
(Indirect, X)	AND (Oper, X)	21	2	6	√	-	-	-	-	-	√	-
(Indirect), Y	AND (Oper), Y	31	2	5*	√	-	-	-	-	-	√	-
(Indirect) <i>NEW</i>	AND (Oper)	32	2	5	√	-	-	-	-	-	√	-

\* => Add 1 if page boundary is crossed

Table (2): AND – Short Reference

Example Immediate (assumed A is \$C5):

Address	Bytes	Mnemonic	
\$9000	29 71	AND #\$71	;
\$9002	...	<i>next OP</i>	; A is now \$41, N=0, Z=0

1100 0101 (\$C5)  
 AND 0111 0001 (\$71)  
0100 0001 (\$41), N=0, Z=0

# ASL

Operation: Shift Left One Bit (Memory or Accumulator) ( $C \leftarrow \boxed{7} \boxed{6} \boxed{5} \boxed{4} \boxed{3} \boxed{2} \boxed{1} \boxed{0} \leftarrow 0$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Accumulator	ASL A	0A	1	2	√	-	-	-	-	-	√	√
Zero Page	ASL Oper	06	2	5	√	-	-	-	-	-	√	√
Zero Page, X	ASL Oper, X	16	2	6	√	-	-	-	-	-	√	√
Absolute	ASL Oper	0E	3	6	√	-	-	-	-	-	√	√
Absolute, X	ASL Oper, X	1E	3	7	√	-	-	-	-	-	√	√

Table (3): ASL – Short Reference

Example Accumulator (assumed A=\$55):

Address	Bytes	Mnemonic	
\$9000	0A	ASL A	;
\$9001	...	<i>next OP</i>	; A is now \$AA, N=1, Z=0

ASL 0101 0101 (\$55)  
1010 1010 (\$AA), C=0, N=1, Z=0

## BBR

Operation: Branch on memory bit x is reset, where x=0-7 MSB of op code (Branch on Mx = 0)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Relative	BBR(0-7) ZP,offs8	0F-7F	3	5*	-	-	-	-	-	-	-	-

\* => Add 1 if branch occurs to same page

\* => Add 2 if branch occurs to different page

Table (4): BBR – Short Reference

Example Relative (assumed \$0032=\$24, same page):

Address	Bytes	Mnemonic	
\$9000	2F 32 10	BBR2 \$32, \$10	; BBR to same page, - NO BRANCH - jumps to \$9003
\$9003	...	<i>next OP</i>	; M2(\$0032) = 1 => no branch, 5 cycles

Example Relative (assumed \$0032=\$24, different page):

Address	Bytes	Mnemonic	
\$90FF	5F 32 10	BBR5 \$32, \$10	; BBR to different page, - NO BRANCH – jumps to \$9102
\$9102	...	<i>next OP</i>	; M5(\$0032) = 1 => no branch, 5 cycles (!!!)

Example Relative (assumed \$0032=\$24, same page):

Address	Bytes	Mnemonic	
\$9000	0F 32 10	BBR0 \$32, \$10	; BBR to same page, - BRANCH - jumps to \$9003 + \$10
...			
\$9013	...	<i>next OP</i>	; M0(\$0032)=0 => branch, 6 cycles

Example Relative (assumed \$0032=\$24, different page):

Address	Bytes	Mnemonic	
\$90FC	0F 32 10	BBR0 \$32, \$10	; BBR to different page, - BRANCH - jumps to \$900F + \$10
...			
\$910F	...	<i>next OP</i>	; M0(\$0032)=0 => branch, 7 cycles

## BBS

Operation: Branch on memory bit x is set, where x=0-7 MSB-8 of op code (Branch on Mx = 1)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Relative	BBS(0-7) ZP,offs8	8F-FF	3	5*	-	-	-	-	-	-	-	-

\* => Add 1 if branch occurs to same page

\* => Add 2 if branch occurs to different page

Table (5): BBS – Short Reference

Example Relative (assumed \$0032=\$24, same page):

Address	Bytes	Mnemonic	
\$9000	8F 32 10	BBS0 \$32, \$10	; BBS to same page, - NO BRANCH - jumps to \$9003
\$9003	...	<i>next OP</i>	; M0(\$0032) =0 => no branch, 5 cycles

Example Relative (assumed \$0032=\$24, different page):

Address	Bytes	Mnemonic	
\$90FF	FF 32 10	BBS7 \$32, \$10	; BBS to different page, - NO BRANCH – jumps to \$9102
\$9102	...	<i>next OP</i>	; M7(\$0032) =0 => no branch, 5 cycles (!!!)

Example Relative (assumed \$0032=\$24, same page):

Address	Bytes	Mnemonic	
\$9000	AF 32 10	BBS2 \$32, \$10	; BBS to same page, - BRANCH - jumps to \$9003 + \$10
...			
\$9013	...	<i>next OP</i>	; M2(\$0032)=1 => branch, 6 cycles

Example Relative (assumed \$0032=\$24, different page):

Address	Bytes	Mnemonic	
\$90FC	DF 32 10	BBS5 \$32, \$10	; BBS to different page, - BRANCH - jumps to \$900F + \$10
...			
\$910F	...	<i>next OP</i>	; M5(\$0032)=1 => branch, 7 cycles

## BCC

Operation: Branch on Carry Clear (Branch on C = 0)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Relative	BCC Oper	90	2	2*	-	-	-	-	-	-	-	-

\* => Add 1 if branch occurs to same page

\* => Add 2 if branch occurs to different page

Table (6): BCC – Short Reference

Example Relative (assumed C=1, no branch):

Address	Bytes	Mnemonic	
\$9000	90 10	BCC \$10	; BCC to same page, - NO BRANCH - jumps to \$9002
\$9002	...	<i>next OP</i>	; C=1 => no branch, 2 cycles

Example Relative (assumed C=1, no branch):

Address	Bytes	Mnemonic	
\$90FF	90 10	BCC \$10	; BCC to same page, - NO BRANCH - jumps to \$9101
\$9101	...	<i>next OP</i>	; C=1 => no branch, 2 cycles

Example Relative (assumed C=0, branch to same page):

Address	Bytes	Mnemonic	
\$9000	90 10	BCC \$10	; BCC to same page, - BRANCH - jumps to \$9002 + \$10
...			
\$9012	...	<i>next OP</i>	; C=0 => branch, 3 cycles

Example Relative (assumed C=0, branch to different page):

Address	Bytes	Mnemonic	
\$90FD	90 10	BCC \$10	; BCC to different page, - BRANCH - jumps to \$9109
...			
\$9109	...	<i>next OP</i>	; C=0 => branch, 4 cycles

## BCS

Operation: Branch on Carry Set (Branch on C = 1)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Relative	BCS Oper	B0	2	2*	-	-	-	-	-	-	-	-

\* => Add 1 if branch occurs to same page

\* => Add 2 if branch occurs to different page

Table (7): BCS – Short Reference

Example Relative (assumed C=0, same page):

Address	Bytes	Mnemonic	
\$9000	B0 10	BCS \$10	; BCS to same page, - NO BRANCH - jumps to \$9002
\$9002	...	<i>next OP</i>	; C=0 => no branch, 2 cycles

Example Relative (assumed C=0, different page):

Address	Bytes	Mnemonic	
\$90FF	B0 10	BCS \$10	; BCS to different page, - NO BRANCH - jumps to \$9101
\$9101	...	<i>next OP</i>	; C=0 => no branch, 2 cycles (!!!)

Example Relative (assumed C=1, same page):

Address	Bytes	Mnemonic	
\$9000	B0 10	BCS \$10	; BCS to same page, - BRANCH - jumps to \$9002 + \$10
...			
\$9012	...	<i>next OP</i>	; C=1 => branch, 3 cycles

Example Relative (assumed C=1, different page):

Address	Bytes	Mnemonic	
\$90FD	B0 10	BCS \$10	; BCS to different page, - BRANCH - jumps to \$910F
...			
\$910F	...	<i>next OP</i>	; C=1 => branch, 4 cycles

## BEQ

Operation: Branch on result zero (Branch on  $Z = 1$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Relative	BEQ Oper	F0	2	2*	-	-	-	-	-	-	-	-

\* => Add 1 if branch occurs to same page

\* => Add 2 if branch occurs to different page

Table (8): BEQ – Short Reference

Example Relative (assumed  $Z=0$ , same page):

Address	Bytes	Mnemonic	
\$9000	F0 10	BEQ \$10	; BEQ to same page, - NO BRANCH - jumps to \$9002
\$9002	...	<i>next OP</i>	; $Z=0$ => no branch, 2 cycles

Example Relative (assumed  $Z=0$ , different page):

Address	Bytes	Mnemonic	
\$90FF	F0 10	BEQ \$10	; BEQ to different page, - NO BRANCH - jumps to \$9101
\$9101	...	<i>next OP</i>	; $Z=0$ => no branch, 2 cycles (!!!)

Example Relative (assumed  $Z=1$ , same page):

Address	Bytes	Mnemonic	
\$9000	F0 10	BEQ \$10	; BEQ to same page, - BRANCH - jumps to \$9002 + \$10
...			
\$9012	...	<i>next OP</i>	; $Z=1$ => branch, 3 cycles

Example Relative (assumed  $Z=1$ , different page):

Address	Bytes	Mnemonic	
\$90FD	F0 10	BEQ \$10	; BEQ to different page, - BRANCH - jumps to \$910F
...			
\$910F	...	<i>next OP</i>	; $Z=1$ => branch, 4 cycles

## BIT

Operation (ZP and ABS): Bit 6 and 7 of the operand are transferred to the status register. If the result of  $A \cap M = 0$  then  $Z = 1$ , otherwise  $Z = 0$  ( $M7 \rightarrow N$ ,  $M6 \rightarrow V$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Immediate <i>NEW</i>	BIT # Oper	89	2	2	-	-	-	-	-	-	√	-
Zero Page	BIT Oper	24	2	3	7	6	-	-	-	-	√	-
Zero Page, X <i>NEW</i>	BIT Oper, X	34	2	4	7	6	-	-	-	-	√	-
Absolute	BIT Oper	2C	3	4	7	6	-	-	-	-	√	-
Absolute, X <i>NEW</i>	BIT Oper, X	3C	3	4*	7	6	-	-	-	-	√	-

\* => Add 1 if page boundary is crossed

Table (9): BIT – Short Reference

**In mode “Immediate” BIT leaves flags N&V unchanged. This is a special case for R65C02!**



## BMI

Operation: Branch on result minus (Branch on N = 1)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Relative	BMI Oper	30	2	2*	-	-	-	-	-	-	-	-

\* => Add 1 if branch occurs to same page

\* => Add 2 if branch occurs to different page

Table (10): BMI – Short Reference

Example Relative (assumed Z=0, same page):

Address	Bytes	Mnemonic	
\$9000	30 10	BMI \$10	; BMI to same page, - NO BRANCH - jumps to \$9002
\$9002	...	<i>next OP</i>	; N=0 => no branch, 2 cycles

Example Relative (assumed Z=0, different page):

Address	Bytes	Mnemonic	
\$90FF	30 10	BMI \$10	; BMI to different page, - NO BRANCH - jumps to \$9101
\$9101	...	<i>next OP</i>	; N=0 => no branch, 2 cycles (!!!)

Example Relative (assumed Z=1, same page):

Address	Bytes	Mnemonic	
\$9000	30 10	BMI \$10	; BMI to same page, - BRANCH - jumps to \$9002 + \$10
...			
\$9012	...	<i>next OP</i>	; N=1 => branch, 3 cycles

Example Relative (assumed Z=1, different page):

Address	Bytes	Mnemonic	
\$90FD	30 10	BMI \$10	; BMI to different page, - BRANCH - jumps to \$910F
...			
\$910F	...	<i>next OP</i>	; N=1 => branch, 4 cycles

## BNE

Operation: Branch on result not zero (Branch on  $Z = 0$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Relative	BNE Oper	D0	2	2*	-	-	-	-	-	-	-	-

\* => Add 1 if branch occurs to same page

\* => Add 2 if branch occurs to different page

Table (11): BNE – Short Reference

Example Relative (assumed  $Z=1$ , same page):

Address	Bytes	Mnemonic	
\$9000	D0 10	BNE \$10	; BNE to same page, - NO BRANCH - jumps to \$9002
\$9002	...	<i>next OP</i>	; $Z=1$ => no branch, 2 cycles

Example Relative (assumed  $Z=1$ , different page):

Address	Bytes	Mnemonic	
\$90FF	D0 10	BNE \$10	; BNE to different page, - NO BRANCH - jumps to \$9101
\$9101	...	<i>next OP</i>	; $Z=1$ => no branch, 2 cycles (!!!)

Example Relative (assumed  $Z=0$ , same page):

Address	Bytes	Mnemonic	
\$9000	D0 10	BNE \$10	; BNE to same page, - BRANCH - jumps to \$9002 + \$10
...			
\$9012	...	<i>next OP</i>	; $Z=0$ => branch, 3 cycles

Example Relative (assumed  $Z=0$ , different page):

Address	Bytes	Mnemonic	
\$90FD	D0 10	BNE \$10	; BNE to different page, - BRANCH - jumps to \$910F
...			
\$910F	...	<i>next OP</i>	; $Z=0$ => branch, 4 cycles

## BPL

Operation: Branch on result plus (Branch on N = 0)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Relative	BPL Oper	10	2	2*	-	-	-	-	-	-	-	-

\* => Add 1 if branch occurs to same page

\* => Add 2 if branch occurs to different page

Table (12): BPL – Short Reference

Example Relative (assumed N=1, same page):

Address	Bytes	Mnemonic	
\$9000	10 10	BPL \$10	; BPL to same page, - NO BRANCH - jumps to \$9002
\$9002	...	<i>next OP</i>	; N=1 => no branch, 2 cycles

Example Relative (assumed N=1, different page):

Address	Bytes	Mnemonic	
\$90FF	10 10	BPL \$10	; BPL to different page, - NO BRANCH - jumps to \$9101
\$9101	...	<i>next OP</i>	; N=1 => no branch, 2 cycles (!!!)

Example Relative (assumed N=0, same page):

Address	Bytes	Mnemonic	
\$9000	10 10	BPL \$10	; BPL to same page, - BRANCH - jumps to \$9002 + \$10
...			
\$9012	...	<i>next OP</i>	; N=0 => branch, 3 cycles

Example Relative (assumed N=0, different page):

Address	Bytes	Mnemonic	
\$90FD	10 10	BPL \$10	; BPL to different page, - BRANCH - jumps to \$910F
...			
\$910F	...	<i>next OP</i>	; N=0 => branch, 4 cycles

## BRA

Operation: Branch always (Branch always)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Relative <i>NEW</i>	BRA Oper	80	2	3*	-	-	-	-	-	-	-	-

\* => Add 1 if branch occurs to different page

Table (13): BRA – Short Reference

## BRK

Operation: Forced Interrupt (PC + 2 ↓, P ↓ NV11DIZC, E & B are written as 1 to stack)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	BRK	00	1	7	-	-	-	-	0	1	-	-

Table (14): BRK – Short Reference

## BVC

Operation: Branch on no overflow (Branch on V = 0)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Relative	BVC Oper	50	2	2*	-	-	-	-	-	-	-	-

\* => Add 1 if branch occurs to same page

\* => Add 2 if branch occurs to different page

Table (15): BVC – Short Reference

Example Relative (assumed V=1, same page):

Address	Bytes	Mnemonic	
\$9000	50 10	BVC \$10	; BVC to same page, - NO BRANCH - jumps to \$9002
\$9002	...	<i>next OP</i>	; V=1 => no branch, 2 cycles

Example Relative (assumed V=1, different page):

Address	Bytes	Mnemonic	
\$90FF	50 10	BVC \$10	; BVC to different page, - NO BRANCH - jumps to \$9101
\$9101	...	<i>next OP</i>	; V=1 => no branch, 2 cycles (!!!)

Example Relative (assumed V=0, same page):

Address	Bytes	Mnemonic	
\$9000	50 10	BVC \$10	; BVC to same page, - BRANCH - jumps to \$9002 + \$10
...			
\$9012	...	<i>next OP</i>	; V=0 => branch, 3 cycles

Example Relative (assumed V=0, different page):

Address	Bytes	Mnemonic	
\$90FD	50 10	BVC \$10	; BVC to different page, - BRANCH - jumps to \$910F
...			
\$910F	...	<i>next OP</i>	; V=0 => branch, 4 cycles

## BVS

Operation: Branch on overflow (Branch on V = 1)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Relative	BVS Oper	70	2	2*	-	-	-	-	-	-	-	-

\* => Add 1 if branch occurs to same page

\* => Add 2 if branch occurs to different page

Table (16): BVS – Short Reference

Example Relative (assumed V=0, same page):

Address	Bytes	Mnemonic	
\$9000	70 10	BVS \$10	; BVS to same page, - NO BRANCH - jumps to \$9002
\$9002	...	<i>next OP</i>	; V=0 => no branch, 2 cycles

Example Relative (assumed V=0, different page):

Address	Bytes	Mnemonic	
\$90FF	70 10	BVS \$10	; BVS to different page, - NO BRANCH - jumps to ; \$9101
\$9101	...	<i>next OP</i>	; V=0 => no branch, 2 cycles (!!!)

Example Relative (assumed V=1, same page):

Address	Bytes	Mnemonic	
\$9000	70 10	BVS \$10	; BVS to same page, - BRANCH - jumps to \$9002 + \$10
...			
\$9012	...	<i>next OP</i>	; V=1 => branch, 3 cycles

Example Relative (assumed V=1, different page):

Address	Bytes	Mnemonic	
\$90FD	70 10	BVS \$10	; BVS to different page, - BRANCH - jumps to ; \$910F
...			
\$910F	...	<i>next OP</i>	; V=1 => branch, 4 cycles

## CLC

Operation: Clear Carry flag (0 → C)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	CLC	18	1	2	-	-	-	-	-	-	-	0

Table (17): CLC – Short Reference

Example:

Address	Bytes	Mnemonic	
\$9000	18	CLC	;
\$9001	...	<i>next OP</i>	; C is now 0

## CLD

Operation: Clear Decimal flag (0 → D)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	CLD	D8	1	2	-	-	-	-	0	-	-	-

Table (18): CLD – Short Reference

Example:

Address	Bytes	Mnemonic	
\$9000	D8	CLD	;
\$9001	...	<i>next OP</i>	; D is now 0

## CLI

Operation: Clear Interrupt Disable flag (0 → I)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	CLI	58	1	2	-	-	-	-	-	0	-	-

Table (19): CLI – Short Reference

Example:

<u>Address</u>	<u>Bytes</u>	<u>Mnemonic</u>	
\$9000	58	CLI	;
\$9001	...	<i>next OP</i>	; I is now 0

## CLV

Operation: Clear Overflow flag (0 → O)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	CLV	B8	1	2	-	0	-	-	-	-	-	-

Table (20): CLV – Short Reference

Example:

<u>Address</u>	<u>Bytes</u>	<u>Mnemonic</u>	
\$9000	B8	CLV	;
\$9001	...	<i>next OP</i>	; V is now 0



## CMP

Operation: Compare Memory and Accumulator (A - M)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Immediate	CMP # Oper	C9	2	2	√	-	-	-	-	-	√	√
Zero Page	CMP Oper	C5	2	3	√	-	-	-	-	-	√	√
Zero Page, X	CMP Oper, X	D5	2	4	√	-	-	-	-	-	√	√
Absolute	CMP Oper	CD	3	4	√	-	-	-	-	-	√	√
Absolute, X	CMP Oper, X	DD	3	4*	√	-	-	-	-	-	√	√
Absolute, Y	CMP Oper, Y	D9	3	4*	√	-	-	-	-	-	√	√
(Indirect, X)	CMP (Oper, X)	C1	2	6	√	-	-	-	-	-	√	√
(Indirect), Y	CMP (Oper), Y	D1	2	5*	√	-	-	-	-	-	√	√
(Indirect) <i>NEW</i>	CMP (Oper)	D2	2	5	√	-	-	-	-	-	√	√

\* => Add 1 if page boundary is crossed

Table (21): CMP – Short Reference

## CPX

Operation: Compare Memory and Index X (X - M)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Immediate	CPX # Oper	E0	2	2	√	-	-	-	-	-	√	√
Zero Page	CPX Oper	E4	2	3	√	-	-	-	-	-	√	√
Absolute	CPX Oper	EC	3	4	√	-	-	-	-	-	√	√

Table (22): CPX – Short Reference

## CPY

Operation: Compare Memory and Index Y (Y - M)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Immediate	CPY # Oper	C0	2	2	√	-	-	-	-	-	√	√
Zero Page	CPY Oper	C4	2	3	√	-	-	-	-	-	√	√
Absolute	CPY Oper	CC	3	4	√	-	-	-	-	-	√	√

Table (23): CPY – Short Reference

## DEC

Operation: Decrement Memory by one (M - 1 → M or A - 1 → A)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Accumulator <i>NEW</i>	DEC A	3A	1	2	√	-	-	-	-	-	√	-
Zero Page	DEC Oper	C6	2	5	√	-	-	-	-	-	√	-
Zero Page, X	DEC Oper, X	D6	2	6	√	-	-	-	-	-	√	-
Absolute	DEC Oper	CE	3	6	√	-	-	-	-	-	√	-
Absolute, X	DEC Oper, X	DE	3	7	√	-	-	-	-	-	√	-

Table (24): DEC – Short Reference

## DEX

Operation: Decrement index X by one ( $X - 1 \rightarrow X$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	DEX	CA	1	2	√	-	-	-	-	-	√	-

Table (25): DEX – Short Reference

Example Implied (assume  $X = \$01$ ):

Address	Bytes	Mnemonic	
\$9000	CA	DEX	;
\$9001	...	<i>next OP</i>	; X is now \$00, N=0, Z=1

DEX 0000 0001 (\$01)  
0000 0000 (\$00), N=0, Z=1

## DEY

Operation: Decrement index Y by one ( $Y - 1 \rightarrow Y$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	DEY	88	1	2	√	-	-	-	-	-	√	-

Table (26): DEX – Short Reference

Example Implied (assume  $Y = \$00$ ):

Address	Bytes	Mnemonic	
\$9000	88	DEY	;
\$9001	...	<i>next OP</i>	; Y is now \$FF, N=1, Z=0

DEY 0000 0000 (\$00)  
1111 1111 (\$FF), N=1, Z=0

## EOR

Operation: “Exclusive-Or” memory or immediate value with accumulator A ( $A \vee M \rightarrow A$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Immediate	EOR # Oper	49	2	2	√	-	-	-	-	-	√	-
Zero Page	EOR Oper	45	2	3	√	-	-	-	-	-	√	-
Zero Page, X	EOR Oper, X	55	2	4	√	-	-	-	-	-	√	-
Absolute	EOR Oper	4D	3	4	√	-	-	-	-	-	√	-
Absolute, X	EOR Oper, X	5D	3	4*	√	-	-	-	-	-	√	-
Absolute, Y	EOR Oper, Y	59	3	4*	√	-	-	-	-	-	√	-
(Indirect, X)	EOR (Oper, X)	41	2	6	√	-	-	-	-	-	√	-
(Indirect), Y	EOR (Oper), Y	51	2	5*	√	-	-	-	-	-	√	-
(Indirect) <i>NEW</i>	EOR (Oper)	52	2	5	√	-	-	-	-	-	√	-

\* => Add 1 if page boundary is crossed

Table (27): EOR – Short Reference

Example Immediate (assumed A is \$23):

Address	Bytes	Mnemonic	
\$9000	49 63	EOR # \$63	;
\$9002	...	<i>next OP</i>	; A is now \$40, N=0, Z=0

0110 0011 (\$63)  
 EOR 0010 0011 (\$23)  
0100 0000 (\$40), N=0, Z=0

## INC

Operation: Increment Memory by one ( $M + 1 \rightarrow M$  or  $A + 1 \rightarrow A$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Accumulator <i>NEW</i>	INC A	1A	1	2	√	-	-	-	-	-	√	-
Zero Page	INC Oper	E6	2	5	√	-	-	-	-	-	√	-
Zero Page, X	INC Oper, X	F6	2	6	√	-	-	-	-	-	√	-
Absolute	INC Oper	EE	3	6	√	-	-	-	-	-	√	-
Absolute, X	INC Oper, X	FE	3	7	√	-	-	-	-	-	√	-

Table (28): INC – Short Reference

## INX

Operation: Increment index X by one ( $X + 1 \rightarrow X$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	INX	E8	1	2	√	-	-	-	-	-	√	-

Table (29): INX – Short Reference

Example Implied (assume  $X = \$0C$ ):

Address	Bytes	Mnemonic	
\$9000	E8	INX	;
\$9001	...	<i>next OP</i>	; X is now \$0D, N=0, Z=0

INX 0000 1100 (\$0C)  
0000 1101 (\$0D), N=0, Z=0

## INY

Operation: Increment index Y by one ( $Y + 1 \rightarrow Y$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	INY	C8	1	2	√	-	-	-	-	-	√	-

Table (30): INY – Short Reference

Example Implied (assume  $Y = \$FF$ ):

Address	Bytes	Mnemonic
\$9000	C8	INY
\$9001	...	<i>next OP</i>

INY 1111 1111 (\$FF)  
0000 0000 (\$00), N=0, Z=1

## JMP

Operation: Jump to new location ( $(PC + 1) \rightarrow PCL$ ,  $(PC + 2) \rightarrow PCH$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Absolute	JMP Oper	4C	3	3	-	-	-	-	-	-	-	-
(Absolute, X) <i>NEW</i>	JMP (Oper, X)	7C	3	6	-	-	-	-	-	-	-	-
Indirect	JMP (Oper)	6C	3	6	-	-	-	-	-	-	-	-

Table (31): JMP – Short Reference

In mode “Indirect” JMP fetches the high-order byte of the effective address from the first byte of the next page, if the indirect address location spans a page boundary. The 6502 fetches it from the first byte of the current page.

R65C02: JMP (\$02FF) gets ADL from \$02FF, ADH from \$0300 (6 cycles)

6502: JMP (\$02FF) gets ADL from \$02FF, ADH from \$0200 (5 cycles)

## JSR

Operation: Jump to subroutine saving return address ( $PC + 2 \downarrow$ ,  $(PC + 1) \rightarrow PCH$ ,  $(PC + 2) \rightarrow PCL$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Absolute	JSR Oper	20	3	6	-	-	-	-	-	-	-	-

Table (32): JSR – Short Reference

Example:

Address	Bytes	Mnemonic
\$9000	20 72 F0	JSR \$F072 ;
F072	...	next OP

## LDA

Operation: Load accumulator A with memory or immediate value ( $M \rightarrow A$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Immediate	LDA # Oper	A9	2	2	√	-	-	-	-	-	√	-
Zero Page	LDA Oper	A5	2	3	√	-	-	-	-	-	√	-
Zero Page, X	LDA Oper, X	B5	2	4	√	-	-	-	-	-	√	-
Absolute	LDA Oper	AD	3	4	√	-	-	-	-	-	√	-
Absolute, X	LDA Oper, X	BD	3	4*	√	-	-	-	-	-	√	-
Absolute, Y	LDA Oper, Y	B9	3	4*	√	-	-	-	-	-	√	-
(Indirect, X)	LDA (Oper, X)	A1	2	6	√	-	-	-	-	-	√	-
(Indirect), Y	LDA (Oper), Y	B1	2	5*	√	-	-	-	-	-	√	-
(Indirect) <i>NEW</i>	LDA (Oper)	B2	2	5	√	-	-	-	-	-	√	-

\* => Add 1 if page boundary is crossed

Table (33): LDA – Short Reference

Example Immediate:

Address	Bytes	Mnemonic
\$9000	A9 00	LDA #\$00 ;
\$9002	...	next OP ; A is now \$00, N=0, Z=1

## LDX

Operation: Load index X with memory or immediate value ( $M \rightarrow X$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Immediate	LDX # Oper	A2	2	2	√	-	-	-	-	-	√	-
Zero Page	LDX Oper	A6	2	3	√	-	-	-	-	-	√	-
Zero Page, Y	LDX Oper, Y	B6	2	4	√	-	-	-	-	-	√	-
Absolute	LDX Oper	AE	3	4	√	-	-	-	-	-	√	-
Absolute, Y	LDX Oper, Y	BE	3	4*	√	-	-	-	-	-	√	-

\* => Add 1 if page boundary is crossed

Table (34): LDX – Short Reference

Example Immediate:

Address	Bytes	Mnemonic	
\$9000	A2 8F	LDX #8F	;
\$9002	...	next OP	; X is now 8F, N=1, Z=0

## LDY

Operation: Load index Y with memory or immediate value ( $M \rightarrow Y$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Immediate	LDY # Oper	A0	2	2	√	-	-	-	-	-	√	-
Zero Page	LDY Oper	A4	2	3	√	-	-	-	-	-	√	-
Zero Page, X	LDY Oper, X	B4	2	4	√	-	-	-	-	-	√	-
Absolute	LDY Oper	AC	3	4	√	-	-	-	-	-	√	-
Absolute, X	LDY Oper, X	BC	3	4*	√	-	-	-	-	-	√	-

\* => Add 1 if page boundary is crossed

Table (35): LDY – Short Reference

Example Immediate:

Address	Bytes	Mnemonic	
\$9000	A0 02	LDY #02	;
\$9002	...	next OP	; Y is now 02, N=0, Z=0



## LSR

Operation: Shift Right One Bit (Memory or Accumulator) (0 → 

7	6	5	4	3	2	1	0
---	---	---	---	---	---	---	---

 → C)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Accumulator	LSR A	4A	1	2	0	-	-	-	-	-	√	√
Zero Page	LSR Oper	46	2	5	0	-	-	-	-	-	√	√
Zero Page, X	LSR Oper, X	56	2	6	0	-	-	-	-	-	√	√
Absolute	LSR Oper	4E	3	6	0	-	-	-	-	-	√	√
Absolute, X	LSR Oper, X	5E	3	7	0	-	-	-	-	-	√	√

Table (36): LSR – Short Reference

## NOP

Operation: No Operation

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	NOP	x3 xB	1	1	-	-	-	-	-	-	-	-
Implied	NOP	EA	1	2	-	-	-	-	-	-	-	-
Implied	NOP	02 22 42 62 82 C2 E2	1(+1)	2	-	-	-	-	-	-	-	-
Implied	NOP	44	1(+1)	3	-	-	-	-	-	-	-	-
Implied	NOP	54 D4 F4	1(+1)	4	-	-	-	-	-	-	-	-
Implied	NOP	DC FC	1(+2)	4	-	-	-	-	-	-	-	-
Implied	NOP	5C	1(+2)	8	-	-	-	-	-	-	-	-

Table (37): NOP – Short Reference (“+1” or “+2” add 1 or 2 to the PC for fetching the next instruction)

## ORA

Operation: “Or” memory or immediate value with accumulator A ( $A \vee M \rightarrow A$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Immediate	OR # Oper	09	2	2	√	-	-	-	-	-	√	-
Zero Page	OR Oper	05	2	3	√	-	-	-	-	-	√	-
Zero Page, X	OR Oper, X	15	2	4	√	-	-	-	-	-	√	-
Absolute	OR Oper	0D	3	4	√	-	-	-	-	-	√	-
Absolute, X	OR Oper, X	1D	3	4*	√	-	-	-	-	-	√	-
Absolute, Y	OR Oper, Y	19	3	4*	√	-	-	-	-	-	√	-
(Indirect, X)	OR (Oper, X)	01	2	6	√	-	-	-	-	-	√	-
(Indirect), Y	OR (Oper), Y	11	2	5*	√	-	-	-	-	-	√	-
(Indirect) <i>NEW</i>	OR (Oper)	12	2	5	√	-	-	-	-	-	√	-

\* => Add 1 if page boundary is crossed

Table (38): OR – Short Reference

Example Immediate (assumed A is \$8A):

Address	Bytes	Mnemonic	
\$9000	09 11	OR #\$11	;
\$9002	...	<i>next OP</i>	; A is now \$9B, N=1, Z=0

1000 1010 (\$8A)  
 OR 0001 0001 (\$11)  
1001 1011 (\$9B), N=1, Z=0

## PHA

Operation: Push accumulator on stack ( $A \downarrow$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	PHA	48	1	3	-	-	-	-	-	-	-	-

Table (39): PHA – Short Reference

## PHP

Operation: Push processor status on stack (P ↓ NV11DIZC, E & B are written as 1 to stack)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	PHP	08	1	3	-	-	-	-	-	-	-	-

Table (40): PHP – Short Reference

## PHX

Operation: Push index register X on stack (X ↓)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied <i>NEW</i>	PHX	DA	1	3	-	-	-	-	-	-	-	-

Table (41): PHX – Short Reference

## PHY

Operation: Push index register Y on stack (Y ↓)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied <i>NEW</i>	PHY	5A	1	3	-	-	-	-	-	-	-	-

Table (42): PHY – Short Reference

## PLA

Operation: Pull accumulator from stack (A ↑)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	PLA	68	1	4	√	-	-	-	-	-	√	-

Table (43): PLA – Short Reference

## PLP

Operation: Pull processor status from stack (P ↑)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	PLP	28	1	4	From Stack							

Table (44): PLP – Short Reference

## PLX

Operation: Pull index register X from stack (X ↑)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied <i>NEW</i>	PLX	FA	1	4	√	-	-	-	-	-	√	-

Table (45): PLX – Short Reference

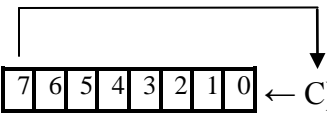
## PLY

Operation: Pull index register Y from stack (Y ↑)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied <i>NEW</i>	PLY	7A	1	4	√	-	-	-	-	-	√	-

Table (46): PLY – Short Reference

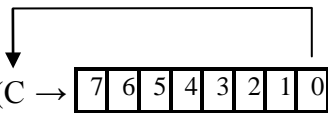
# ROL

Operation: Rotate one bit left (memory or accumulator) ()

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Accumulator	ROL A	2A	1	2	√	-	-	-	-	-	√	√
Zero Page	ROL Oper	26	2	5	√	-	-	-	-	-	√	√
Zero Page, X	ROL Oper, X	36	2	6	√	-	-	-	-	-	√	√
Absolute	ROL Oper	2E	3	6	√	-	-	-	-	-	√	√
Absolute, X	ROL Oper, X	3E	3	7	√	-	-	-	-	-	√	√

Table (47): ROL – Short Reference

# ROR

Operation: Rotate one bit right (memory or accumulator) ()

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Accumulator	ROR A	6A	1	2	√	-	-	-	-	-	√	√
Zero Page	ROR Oper	66	2	5	√	-	-	-	-	-	√	√
Zero Page, X	ROR Oper, X	76	2	6	√	-	-	-	-	-	√	√
Absolute	ROR Oper	6E	3	6	√	-	-	-	-	-	√	√
Absolute, X	ROR Oper, X	7E	3	7	√	-	-	-	-	-	√	√

Table (48): ROR – Short Reference

## RMB

Operation: Reset memory bit x. Where x = 0-7 MSB of op code (0 → Mx)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Zero Page <i>NEW</i>	RMB(0-7) Oper	07-77	2	5	-	-	-	-	-	-	-	-

Table (49): RMB – Short Reference

Example Zero Page:

Address	Bytes	Mnemonic	
\$9000	37 14	RMB3 \$14	; \$0014 = \$C9
\$9002	...	<i>next OP</i>	; \$0014 is now \$C1

## RTI

Operation: Return from interrupt (P ↑, PC ↑)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	RTI	40	1	6	From Stack							

Table (50): RTI – Short Reference

## RTS

Operation: Return from subroutine (PC ↑, PC + 1 → PC)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	RTS	60	1	6	-	-	-	-	-	-	-	-

Table (51): RTS – Short Reference

## SBC

Operation: Subtract memory or immediate value from accumulator with borrow ( $A - M - \bar{C} \rightarrow A, C$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Immediate	SBC # Oper	E9	2	2+	√	√	-	-	-	-	√	√
Zero Page	SBC Oper	E5	2	3+	√	√	-	-	-	-	√	√
Zero Page, X	SBC Oper, X	F5	2	4+	√	√	-	-	-	-	√	√
Absolute	SBC Oper	ED	3	4+	√	√	-	-	-	-	√	√
Absolute, X	SBC Oper, X	FD	3	4+*	√	√	-	-	-	-	√	√
Absolute, Y	SBC Oper, Y	F9	3	4+*	√	√	-	-	-	-	√	√
(Indirect, X)	SBC (Oper, X)	E1	2	6+	√	√	-	-	-	-	√	√
(Indirect), Y	SBC (Oper), Y	F1	2	5+*	√	√	-	-	-	-	√	√
(Indirect) <i>NEW</i>	SBC (Oper)	F2	2	5+	√	√	-	-	-	-	√	√

\* => Add 1 if page boundary is crossed

+ => Add 1 if mode is "DECIMAL" (D='1')

Table (52): SBC – Short Reference

## SEC

Operation: Set Carry flag ( $1 \rightarrow C$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	SEC	38	1	2	-	-	-	-	-	-	-	1

Table (53): SEC – Short Reference

Example:

Address	Bytes	Mnemonic	
\$9000	38	SEC	;
\$9001	...	<i>next OP</i>	; C is now 1

## SED

Operation: Set Decimal flag (1 → D)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	SED	F8	1	2	-	-	-	-	1	-	-	-

Table (54): CLD – Short Reference

Example:

<u>Address</u>	<u>Bytes</u>	<u>Mnemonic</u>	
\$9000	F8	SED	;
\$9001	...	<i>next OP</i>	; D is now 1

## SEI

Operation: Set Interrupt Disable flag (1 → I)

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	SEI	78	1	2	-	-	-	-	-	1	-	-

Table (55): SEI – Short Reference

Example:

<u>Address</u>	<u>Bytes</u>	<u>Mnemonic</u>	
\$9000	78	SEI	;
\$9001	...	<i>next OP</i>	; I is now 1



## SMB

Operation: Set memory bit  $x$ . Where  $x = 0-7$  (MSB – 8) of op code ( $1 \rightarrow Mx$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Zero Page <i>NEW</i>	SMB(0-7) Oper	87-F7	2	5	-	-	-	-	-	-	-	-

Table (56): SMB – Short Reference

Example Zero Page:

Address	Bytes	Mnemonic	
\$9000	F7 14	SMB \$F0	; \$00F0 = \$43
\$9002	...	<i>next OP</i>	; \$00F0 is now \$C3

## STA

Operation: Store accumulator in memory ( $A \rightarrow M$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Zero Page	STA Oper	85	2	3	-	-	-	-	-	-	-	-
Zero Page, X	STA Oper, X	95	2	4	-	-	-	-	-	-	-	-
Absolute	STA Oper	8D	3	4	-	-	-	-	-	-	-	-
Absolute, X	STA Oper, X	9D	3	5	-	-	-	-	-	-	-	-
Absolute, Y	STA Oper, Y	99	3	5	-	-	-	-	-	-	-	-
(Indirect, X)	STA (Oper, X)	81	2	6	-	-	-	-	-	-	-	-
(Indirect), Y	STA (Oper), Y	91	2	6	-	-	-	-	-	-	-	-
(Indirect) <i>NEW</i>	STA (Oper)	92	2	5	-	-	-	-	-	-	-	-

Table (57): STA – Short Reference

## STX

Operation: Store index X in memory ( $X \rightarrow M$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Zero Page	STX Oper	86	2	3	-	-	-	-	-	-	-	-
Zero Page, Y	STX Oper, Y	96	2	4	-	-	-	-	-	-	-	-
Absolute	STX Oper	8E	3	4	-	-	-	-	-	-	-	-

Table (58): STX – Short Reference

## STY

Operation: Store index Y in memory ( $Y \rightarrow M$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Zero Page	STY Oper	84	2	3	-	-	-	-	-	-	-	-
Zero Page, X	STY Oper, X	94	2	4	-	-	-	-	-	-	-	-
Absolute	STY Oper	8C	3	4	-	-	-	-	-	-	-	-

Table (59): STY – Short Reference

## STZ

Operation: Store value # $\$00$  in memory ( $\#\$00 \rightarrow M$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Zero Page <i>NEW</i>	STZ Oper	64	2	3	-	-	-	-	-	-	-	-
Zero Page, X <i>NEW</i>	STZ Oper, X	74	2	4	-	-	-	-	-	-	-	-
Absolute <i>NEW</i>	STZ Oper	9C	3	4	-	-	-	-	-	-	-	-
Absolute, X <i>NEW</i>	STZ Oper, X	9E	3	5	-	-	-	-	-	-	-	-

Table (60): STZ – Short Reference

## TAX

Operation: Transfer accumulator A to index X ( $A \rightarrow X$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	TAX	AA	1	2	√	-	-	-	-	-	√	-

Table (61): TAX – Short Reference

Example Implied (assume A=\$5D):

Address	Bytes	Mnemonic	
\$9000	AA	TAX	;
\$9001	...	<i>next OP</i>	; X is now \$5D, N=0, Z=0

## TAY

Operation: Transfer accumulator A to index Y ( $A \rightarrow Y$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	TAY	A8	1	2	√	-	-	-	-	-	√	-

Table (62): TAY – Short Reference

Example Implied (assume A=\$89):

Address	Bytes	Mnemonic	
\$9000	A8	TAY	;
\$9001	...	<i>next OP</i>	; Y is now \$89, N=1, Z=0

## TRB

Operation: Test memory with accumulator A and reset memory bits. If the result of  $A \cap M = 0$  then  $Z = 1$ , otherwise  $Z = 0$  ( $\overline{A} \cap M \rightarrow M$ )

Addressing Mode		Assembly Language Form		OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Zero Page	<i>NEW</i>	TRB	Oper	14	2	5	-	-	-	-	-	-	√	-
Absolute	<i>NEW</i>	TRB	Oper	1C	3	6	-	-	-	-	-	-	√	-

Table (63): TRB – Short Reference

Example Zero Page (assume A=\$03):

<u>Address</u>	<u>Bytes</u>	<u>Mnemonic</u>
\$9000	14 D2	TRB \$D2 ; \$00D2 = \$61
\$9002	...	<i>next OP</i> ; \$00D2 is now \$60, Z=0

```

0000 0011 (A=$03)
NOT 1111 1100
AND 0110 0001 (M=$61)
0110 0000 ($60) New value for memory
    
```

```

0000 0011 (A=$03)
AND 0110 0001 (M=$61)
0000 0001 Z=0 New value for flag Z
    
```

## TSB

Operation: Test memory with accumulator A and set memory bits. If the result of  $A \cap M = 0$  then  $Z = 1$ , otherwise  $Z = 0$  ( $A \cup M \rightarrow M$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Zero Page <i>NEW</i>	TSB Oper	04	2	5	-	-	-	-	-	-	√	-
Absolute <i>NEW</i>	TSB Oper	0C	3	6	-	-	-	-	-	-	√	-

Table (64): TSB – Short Reference

Example Zero Page (assume A=\$0C):

Address	Bytes	Mnemonic
\$9000	04 D2	TSB \$D2 ; \$00D2 = \$11
\$9002	...	<i>next OP</i> ; \$00D2 is now \$1D, Z=1

0000 1100 (A=\$0C)  
OR 0001 0001 (M=\$11)  
0001 1101 (\$1D) New value for memory

0000 1100 (A=\$0C)  
AND 0001 0001 (M=\$11)  
0000 0000 Z=1 New value for flag Z

## TSX

Operation: Transfer stack pointer SP to index X ( $SP \rightarrow X$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	TSX	BA	1	2	√	-	-	-	-	-	√	-

Table (65): TSX – Short Reference

Example Implied (assume SP=\$F2):

Address	Bytes	Mnemonic
\$9000	BA	TSX
\$9001	...	<i>next OP</i> ; X is now \$F2, N=1, Z=0

## TXA

Operation: Transfer index X to accumulator A ( $X \rightarrow A$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	TXA	8A	1	2	√	-	-	-	-	-	√	-

Table (66): TXA – Short Reference

Example Implied (assume X=\$00):

Address	Bytes	Mnemonic	
\$9000	8A	TXA	
\$9001	...	<i>next OP</i>	; A is now \$00, N=0, Z=1

## TXS

Operation: Transfer index X to stack pointer SP ( $X \rightarrow SP$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	TXS	9A	1	2	-	-	-	-	-	-	-	-

Table (67): TXS – Short Reference

Example Implied (assume X=\$E0):

Address	Bytes	Mnemonic	
\$9000	9A	TXS	
\$9001	...	<i>next OP</i>	; SP is now \$E0 (all flags are uneffected)

## TYA

Operation: Transfer index Y to accumulator A ( $Y \rightarrow A$ )

Addressing Mode	Assembly Language Form	OP CODE	No. Bytes	No. Cycles	N	V	E	B	D	I	Z	C
Implied	TYA	98	1	2	√	-	-	-	-	-	√	-

Table (68): TYA – Short Reference

Example Implied (assume  $Y=\$14$ ):

Address	Bytes	Mnemonic	
\$9000	98	TYA	;
\$9001	...	<i>next OP</i>	; A is now \$14, N=0, Z=0

## 4

# Registers

## List of Registers

Name	Address	Width	Access	Description
A	-	8	Instruction	Accumulator
X	-	8	Instruction	X Register
Y	-	8	Instruction	Y Register
S	-	8	Instruction	Stack
PSW	-	8	Instruction	Processor Status Word
PC	-	16	Instruction	Program Counter

Table 1: List of registers

## PSW – Description

Bit #	Access	Description
7		N-Flag, negative result (1 -> negative)
6		V-Flag, overflow (1 -> overflow)
5		<i>RESERVED (always read as '1')</i>
4		B-Flag, BRK instruction
3		D-Flag, decimal mode (1 -> decimal)
2		I-Flag, interrupt disable (1 -> disabled)
1		Z-Flag, zero result (1 -> zero)
0		C-Flag, carry

Table 2: PSW Description

*Reset Value: 24h*



# 5

## Clocks

Name	Source	Rates (MHz)			Remarks	Description
		Max	Typ	Min		
clk_clk_i	-		110-180	0	Achievable rates depend on synthesis and place&route.	Master Clock

**Table 3: List of clocks**

The following table gives you an overview of standard clocking rates for some different families and vendors (only standalone core – no system):

Vendor	Device	Rates (MHz)			Logic Utilisation	Remarks
		Max	Typ	Min		
Intel	Stratix II: EP2S60F672C5ES		146	0	1161 ALUTs	“SLOW” model
Intel	Cyclone III: EP3C5F256C6		125	0	1593 Elements	“SLOW” model
Intel	MAX V: 5M2210ZF256C4		62	0	1542 Elements	“SLOW” model
Intel	MAX 10: 10M16DCF256C7G		113	0	1553 Elements	“SLOW” model
Intel	Stratix V: 5SGXMA3K2F35C3		242	0	692 ALMs	“SLOW” model
Xilinx	Artix-7: xc7a15tcbg236-3		181	0	1511 LUTs	“Default” strategy
Xilinx	Virtex-7: xc7v585tffg1157-3		144	0	1349 LUTs	“Default” strategy
Xilinx	Kintex-7: xc7k70tcbg484-3		150	0	1349 LUTs	“Default” strategy

**Table 4: List of clock rates**

# 6

## Interrupts

The priority of interrupts is as followed if all of three interrupts occur at the same time:

BRK – NMI – IRQ

In contrast to the chip version of the R65C02, the interrupts are not delay able by some instructions like “CLI”.

### NMI

The interrupt input “nmi\_n\_i” is edge triggered at a falling edge ‘1’ -> ‘0’. It is not mask-able.

There is a need of one cycle ‘1’ following with one cycle ‘0’ to detect a NMI interrupt as “valid”. “nmi\_n\_i” can switch back to ‘1’ or held on ‘0’ within cycle 3.

After the internal IRQ sequence of 7 cycles has finished, the “valid” flag re-set and allow another incoming NMI to detect. The maximum frequency of falling edges on “nmi\_n\_i” must be carefully defined to avoid a stack overflow while a currently processing interrupt is interrupted by another one.

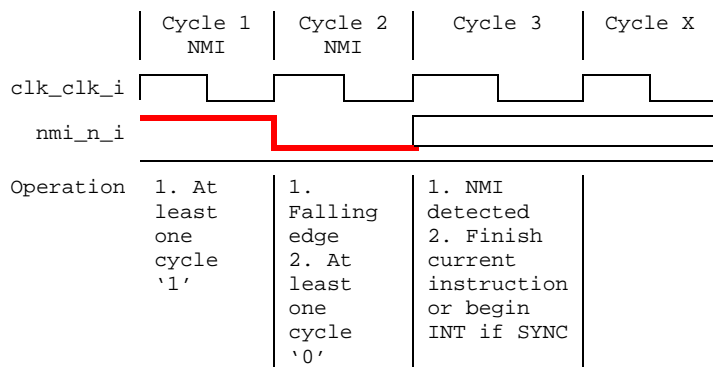


Figure (3): Interrupt NMI – Timing Diagram

## IRQ

The interrupt input “irq\_n\_i” is level triggered at a level of ‘0’. It is mask-able by using the I-Flag (I=’1’ to disable IRQ).

There is a need of level ‘0’ at a rising edge of “clk\_clk\_i” and “SYNC=’1’” to detect an IRQ interrupt as “valid”. “irq\_n\_i” can switch back to ‘1’ or held on ‘0’ within cycle 3.

After the internal IRQ sequence of 7 cycles has finished, the I-Flag is set to ‘1’ automatically to disallow another incoming IRQ to detect. The starting ISR (Interrupt Service Routine) is responsible to manage and service the devices which were generating the interrupt and enable IRQ again (I-Flag=’0’).

Devices must switch back “irq\_n\_i” to ‘1’ after the interrupt service is finished.

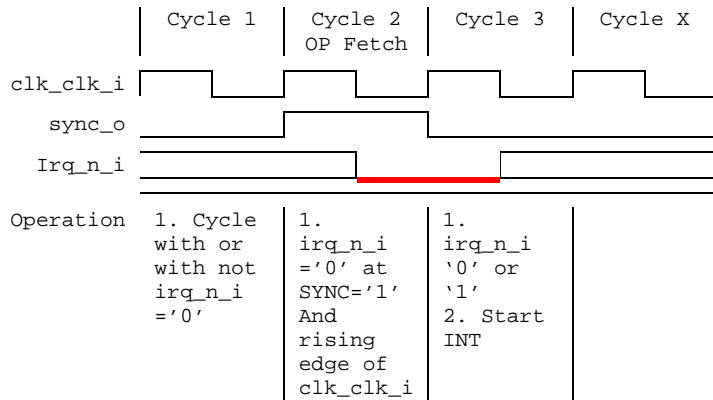


Figure (4): Interrupt IRQ – Timing Diagram

## BRK

“BRK” is an instruction to force an interrupt by software. Internally it is handled like a real hardware interrupt and is mostly used by debug software.

“BRK” transfers the B-Flag = 1 onto the stack. The B-Flag itself leaves untouched and the I-Flag after “BRK” is 1.

# Appendix A

## ADC / SBC Decimal Mode

The R65C02 allows handling of Binary Coded Decimals (BCD) directly if the D-Flag is set to 1.

BCD defines a representation of decimal numbers 0-9 within an underlying hexadecimal number system.

For example:

99d => 63h

But in BCD the representation is

99d => 99h (BCD)

The R65C02 has the ability to convert the result value into BCD automatically if the nibbles exceed the decimal range of 0-9 with some limitations. Input values are not corrected in any way.

Some examples with ADC IMM (DECIMAL):

ACCU	IMM	Carry in	Real World Result HEX	Real World Result BCD	R65C02 Result BCD	Low Nibble (LN) High Nibble (HN)
01h	09h	0	0Ah	10h	10h	LN > 9 (Ah), Carry to HN (=10h)
00h	09h	1	0Ah	10h	10h	LN > 9 (Ah), Carry to HN (=10h)
01h	0Fh	0	10h	16h	00h	LN < 9 (0h), NO Carry to HN (=00h)
0Fh	0Fh	0	1Eh	30h	14h	LN > 9 (Eh), Carry to HN (=1h new), HN <sub>new</sub> = HN <sub>a</sub> + HN <sub>imm</sub> + 1
F0h	F0h	0	(1)E0h	(4)80h	40h	LN < 9 (0h), NO Carry to HN (=Eh new), HN <sub>new</sub> = HN <sub>a</sub> + HN <sub>imm</sub> + 0

FFh	FFh	0	(1)FEh	(5)10h	54h	LN > 9 (Eh), Carry to HN (=Fh new), HN <sub>new</sub> = HN <sub>a</sub> + HN <sub>imm</sub> + 1
FFh	FFh	1	(1)FFh	(5)11h	55h	LN > 9 (Fh), Carry to HN (=Fh new), HN <sub>new</sub> = HN <sub>a</sub> + HN <sub>imm</sub> + 1

**Table 5: List of registers**

In general (NOT valid for 65x02), there are four levels of decoding and corrections for hexadecimal values to convert to BCD:

Level	Result pure HEX		Corrected Result	System	Tasks
L1	04h	=>	04h	(BCD)	no correction needed
L2	0Ch	=>	12h	(BCD)	correction of low nibble, add carry to the high nibble
L3	12h	=>	18h	(BCD)	correction of low nibble from high nibble (backward correction) if carry occur from low to high nibble.
L4	1Bh	=>	27h	(BCD)	correction of low nibble from high nibble , add carry to the high nibble (backward correction)

**Table 6: List of registers**

Even the R65C02 correct the nibbles if a nibble value is greater than 9, the handling and generation of carry bits from one nibble to the other covers not all possible decoding levels L1 to L4. Incorrect results may occur if using invalid BCD numbers containing the hex parts A-F. The backward correction to the low nibble and new generation of the carry to correct the high nibble (L4) was not implemented into the R65C02 chip and the r65c02\_tc IP core also.

# Instruction Table

		LSB															
		0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
0	BRK	Relative	ORA (IND,X)			TSB ZP	ORA ZP	ASL ZP	RMB0 ZP	PHP Implied	ORA IMM	ASL Accum		TSB ABS	ORA	ASL	BBR0 ZP
	1	BPL Relative	ORA (IND,Y)	ORA (IND)		TRB ZP	ORA ZP	ASL ZP	RMB1 ZP	CLC Implied	ORA ABS,Y	INC Accum		TRB ABS	ORA	ASL	BBR1 ZP
2	JSR	Relative	AND (IND,X)			BIT ZP	AND ZP	ROL ZP	RMB2 ZP	PLP Implied	AND IMM	ROL Accum		BIT ABS	AND	ROL	BBR2 ZP
	3	BMI Relative	AND (IND,Y)	AND (IND)		BIT ZP	AND ZP	ROL ZP	RMB3 ZP	SEC Implied	AND ABS,Y	DEC Accum		BIT ABS,X	AND	ROL	BBR3 ZP
4	RTI	Implied	ORA (IND,X)				ORA ZP	LSR ZP	RMB4 ZP	PHA Implied	ORA IMM	LSR Accum		IMP ABS	ORA	LSR	BBR4 ZP
	5	BVC Relative	ORA (IND,Y)	ORA (IND)			ORA ZP	LSR ZP	RMB5 ZP	CLI Implied	ORA ABS,Y	PHY Implied			ORA	LSR	BBR5 ZP
6	RTS	Relative	ADC (IND,X)			STZ ZP	ADC ZP	ROR ZP	RMB6 ZP	PLA Implied	ADC IMM	ROR Accum		JMP (IND)	ADC	ROR	BBR6 ZP
	7	BVS Relative	ADC (IND,Y)	ADC (IND)		STZ ZP	ADC ZP	ROR ZP	RMB7 ZP	SEI Implied	ADC ABS,Y	PLY Implied		JMP (IND,X)	ADC	ROR	BBR7 ZP
8	BRA	Relative	STA (IND,X)			STY ZP	STA ZP	STX ZP	SMB0 ZP	DEY Implied	STA IMM	TXA Implied		STY ABS	STA	STX	BR80 ZP
	9	BCC Relative	STA (IND,Y)	STA (IND)		STY ZP	STA ZP	STX ZP	SMB1 ZP	TYA Implied	STA ABS,Y	TXS Implied		STZ ABS	STA	STZ	BR81 ZP
A	LDY	IMM	LDA (IND,X)	LDA (IND)		LDY ZP	LDA ZP	LDX ZP	SMB2 ZP	TAY Implied	LDA IMM	TAX Implied		LDY ABS	LDA	LDX	BR82 ZP
	B	Relative	LDA (IND,Y)	LDA (IND)		LDY ZP	LDA ZP	LDX ZP	SMB3 ZP	CLV Implied	LDA ABS,Y	TSX Implied		LDY ABS,X	LDA	LDX	BR83 ZP
C	CPY	IMM	CMP (IND,X)			CPY ZP	CMP ZP	DEC ZP	SMB4 ZP	INY Implied	CMP IMM	DEX Implied		CPY ABS	CMP	DEC	BR84 ZP
	D	BNE Relative	CMP (IND,Y)	CMP (IND)		CPY ZP	CMP ZP	DEC ZP	SMB5 ZP	CLD Implied	CMP ABS,Y	PHX Implied		CPY ABS,X	CMP	DEC	BR85 ZP
E	CPX	IMM	SBC (IND,X)			CPX ZP	SBC ZP	INC ZP	SMB6 ZP	INX Implied	SBC IMM	NOP Implied		CPX ABS	SBC	INC	BR86 ZP
	F	BEQ Relative	SBC (IND,Y)	SBC (IND)		CPX ZP	SBC ZP	INC ZP	SMB7 ZP	SED Implied	SBC ABS,Y	PLX Implied		CPX ABS,X	SBC	INC	BR87 ZP
0																	
0																	

+ Add 1 to N if in decimal mode, \* ADD 1 to N if page boundary is crossed  
 \*\* Add 1 to N if branch occurs to same page or add 2 to N if branch occurs to different page

# Index

---

- [1] Rockwell, Product Description R65C02, R65C102 and R65C112 – R65C00 Microprocessors Document No. 29651N52, Rev. 6, June 1987, Order No. 2149
- [2] MOS Technology, MCS6500 Microcomputer Family, Hardware Manual Publication No. 6500-10A, Rev. A, January 1976
- [3] MOS Technology, MCS6500 Microcomputer Family, Software Manual Publication No. 6500-50A, January 1976
- [4] Western Design Center, Programming the 65816 2007